

Lexicographic choice functions without archimedeanity

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Imprecise probabilities

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Basic idea of imprecise probabilities: decisions and choice.

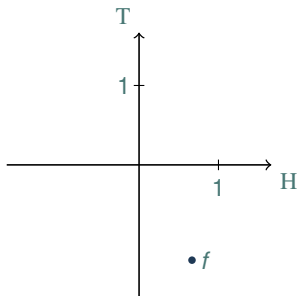
What do we choose between: gambles

The random variable X takes values in the possibility space \mathcal{X} .

A **gamble** $f: \mathcal{X} \rightarrow \mathbb{R}$ is an uncertain reward whose value is $f(X)$, and we collect all gambles in \mathcal{L} .



$$\mathcal{X} = \{H, T\}$$



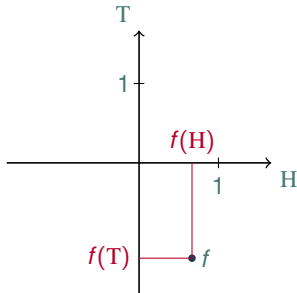
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Which gambles are strictly preferred to 0?

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$$f_1 \succ f_2 \Leftrightarrow f_1 - f_2 \succ 0 \Leftrightarrow f_1 - f_2 \in D$$

for all f_1 and f_2 in \mathcal{L} .

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To summarise:

$$D = \{f \in \mathcal{L} : f \succ 0\}.$$

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Sets of desirable gambles are more general than lower previsions or sets of probabilities.

Working with them is simple and elegant.

They form a strong belief structure: they generalise conservative logical inference (natural extension).

More general choice

A : option set of gambles (non-empty but finite)

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\mathcal{Q} : set of all the non-empty but finite subsets of \mathcal{L} .

A choice function C is a map

$$C: \mathcal{Q} \rightarrow \mathcal{Q} \cup \{\emptyset\}: A \mapsto C(A) \text{ such that } C(A) \subseteq A.$$

Axioms for “coherence”

We call a choice function C on \mathcal{Q} **coherent** if for all A, A_1, A_2 in \mathcal{Q} , f, g in \mathcal{L} and λ in $\mathbb{R}_{>0}$:

C1 $C(A) \neq \emptyset$; [non-emptiness]

C2 if $f < g$ then $\{g\} = C(\{f, g\})$; [non-triviality]

C3a if $A \subseteq R(A_1)$ and $A_1 \subseteq A_2$ then $A \subseteq R(A_2)$; [Sen's condition α]

C3b if $A_1 \subseteq R(A_2)$ and $A \subseteq A_1$
then $A_1 \setminus A \subseteq R(A_2 \setminus A)$; [Aizerman's condition]

C4a if $A_1 \subseteq C(A_2)$ then $\lambda A_1 \subseteq C(\lambda A_2)$; [scaling]

C4b if $A_1 \subseteq C(A_2)$ then $A_1 + \{f\} \subseteq C(A_2 + \{f\})$. [addition]

Coherent choice functions do not form a strong belief structure.

Convexity

We need something stronger: is a **convexity** axiom sufficient?

For all A and A_1 in \mathcal{Q} :

C5 if $A \subseteq A_1 \subseteq \text{CH}(A)$ then $C(A) \subseteq C(A_1)$. [convexity]

Purely binary choice with choice functions

For all A in \mathcal{Q}

$C_D(A) := \{f \in A : (\forall g \in A) g - f \notin D\}$ “the undominated gambles in A ”

But: not every choice function is binary.

Connection between desirability and choice functions

From D to C : $C_D(A) := \{f \in A : (\forall g \in A) g - f \notin D\}$ for all A in \mathcal{Q}

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Coherence is preserved.

Desirability and convexity

If D is coherent, then C_D is coherent.

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C_D is coherent + C5 $\Leftrightarrow D$ is coherent and $\text{posi}(D^c) = D^c$.

Interestingly, $\text{posi}(D^c) = D^c$ is a characterisation of lexicographic probability!

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We collect such D 's in $\bar{\mathcal{D}}_L := \{D \text{ coherent} : \text{posi}(D^c) = D^c\}$, and call $\bar{\mathcal{C}}_L := \{C_D : D \in \bar{\mathcal{D}}_L\}$ the lexicographic choice functions.

Is every coherent and **convex** choice function an infimum of **lexicographic** choice functions?

Rejection sets: choice amongst three

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When $|\mathcal{X}| = 2$, choice functions are determined by choice amongst three options.

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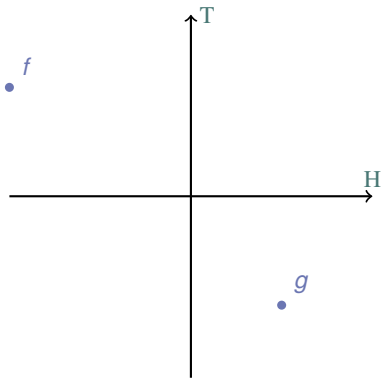
The question becomes: “Is $0 \in R(\{0, f, g\})$?”

Rejection sets: choice amongst three

Is $0 \in R(\{0, f, g\})$? Which gambles allow me to reject 0?

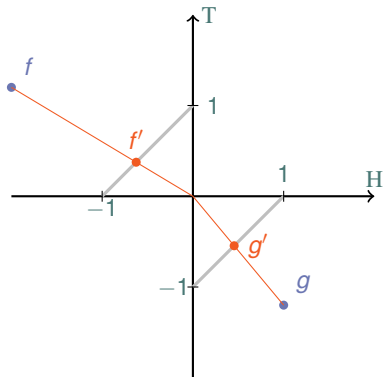
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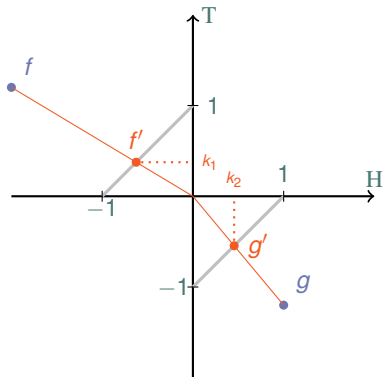
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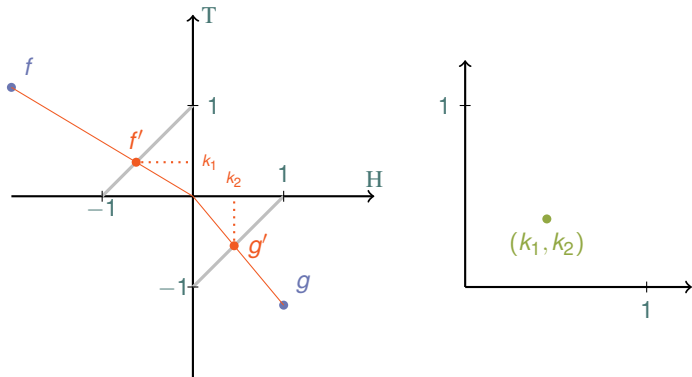
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$$K = \{(k_1, k_2) \in [0, 1]^2 : 0 \in R(\{0, (k_1 - 1, k_1), (k_2, k_2 - 1)\})\}$$

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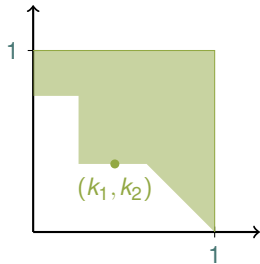
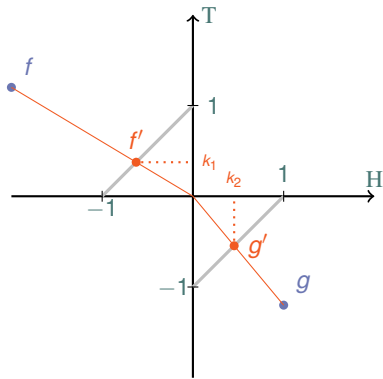
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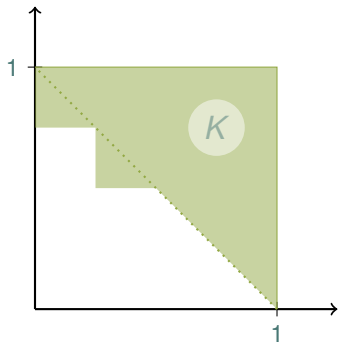
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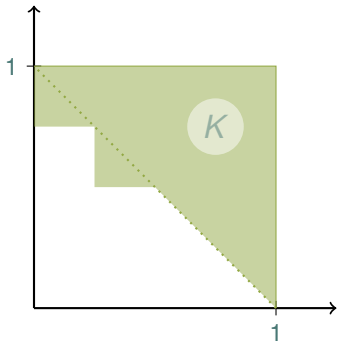


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From coherence axioms to graphical properties



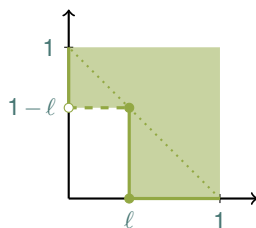
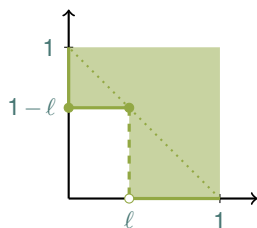
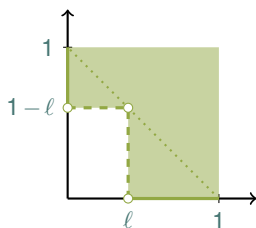
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Properties:

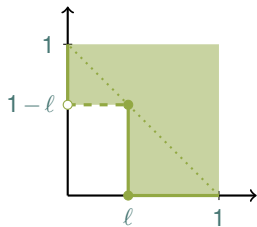
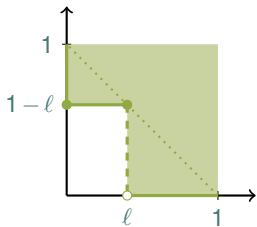
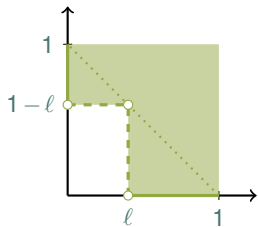
- $(0,0) \notin K$;
- if $(k_1, k_2) \in K$ and $k_1 \leq l_1$ and $k_2 \leq l_2$, then $(l_1, l_2) \in K$ [K is increasing];
- C satisfies C5 $\Leftrightarrow \{(l_1, l_2) \in [0,1]^2 : l_1 + l_2 > 1\} \subseteq K$;
- below the diagonal, the border of K has only horizontal or vertical lines.

What do (infima of) lexicographic choice functions look like?

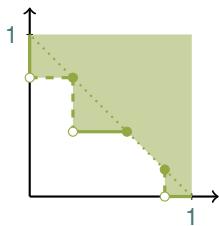
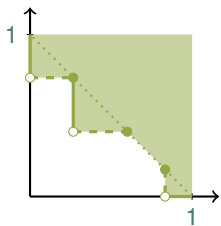
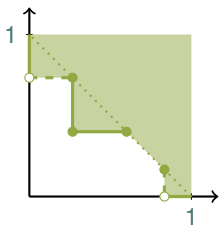


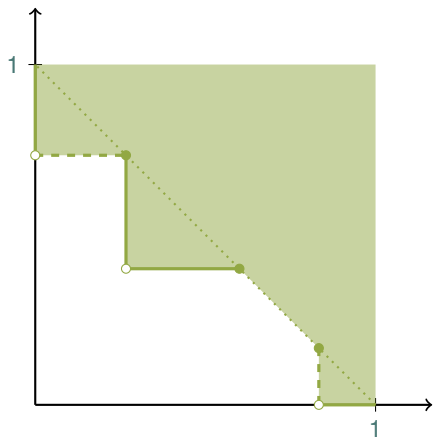
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Conclusion

Coherence and $C5$ is not sufficient to guarantee that C is an infimum of lexicographic choice functions.

