

A queueing model for group-screening facilities

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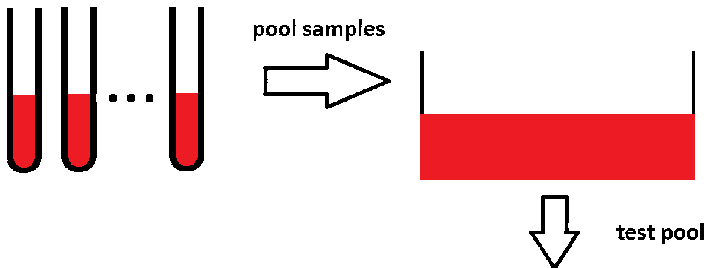


Outline

- 1 Group screening: what?
- 2 Queueing model
- 3 Static versus dynamic group screening
- 4 Comparison of results
 - Testing probability
 - Mean delay
- 5 Conclusion

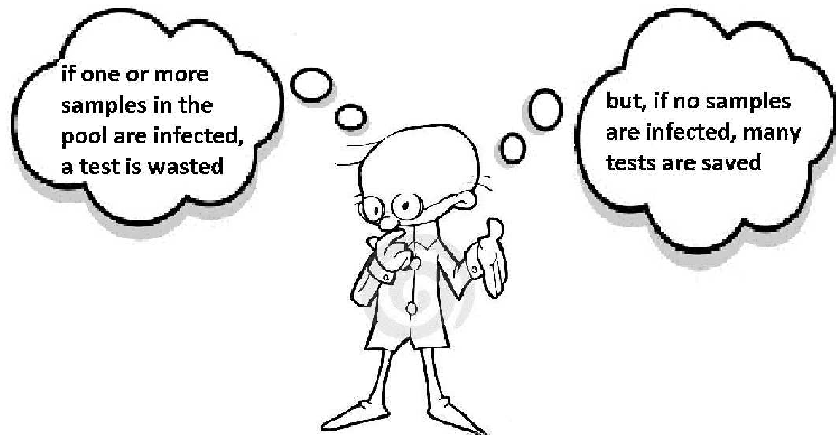
Group screening

- WWII: detect syphilitic men drafted for military service
 - ▶ Expensive
 - ▶ Small prevalence rate
- Idea: test samples **in group**



Infected? No: all samples are not infected
Yes: retest samples (individual/subgroups)

Does it work?



Example

- Prevalence rate 1%
- 5 samples per pool
- Individual retesting



- $\Pr[\text{pool infected}] = 1 - (0.99)^5 \approx 5\%$
- If pool infected: 5 extra tests (+ 1 group test)
- Else: only 1 group test
- $E[\# \text{ tests per pool}] \approx 0.05 \cdot 6 + 0.95 \cdot 1 = 1.25 \text{ instead of } 5$

Other applications of group screening

- Screening for HIV, Influenza, West Nile Virus
- DNA library screening
- Drug discovery
- Quality control

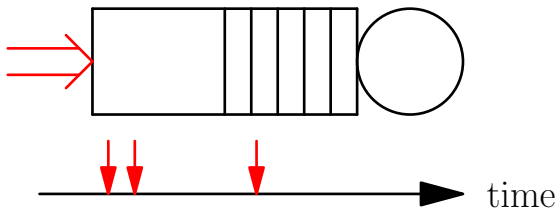
Optimal group size

- Larger group size:
 - ▶ More items (samples) in a group (pool)
 - ▶ Larger probability that group is bad
- Dorfman: standard model
 - ▶ Many items to be screened
 - ▶ Items present from the start
- Practical context usually not static but dynamic
 - ▶ Items arrive over time
 - ▶ Extra decision variable: minimum group size
 - ▶ Queueing model

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- Items arrive spread over time



- Items arrive spread over time
- Await in queue their screening



- Items arrive spread over time
- Await in queue their screening
- By the screening facility



Screening

- In group: batch/bulk service
- Server capacity c : maximum group size
- Minimum batch size l : minimum group size

Screening

- Service time of group: # required tests
- Dependent on # items in the group
- Capture screening policy in distribution screening time

Example: individual retesting

- S_j : service time of a group of j items
- $S_j(z)$: probability generating function of S_j
- $\bar{p} \triangleq 1 - \text{prevalence rate}$
- $S_j(z) = \bar{p}^j z + (1 - \bar{p}^j) z^{j+1}$, $j \geq 2$

Performance measures

- **Testing probability** f : fraction of slots during which server (test unit) is serving (screening)
- **Mean delay** $\bar{D}$: average time between the arrival of an item and the moment at which the result of the item is known

Numerical example: f

	$c = 6$	$c = 7$	$c = 8$	$c = 9$	$c = 10$	$c = 11$	$c = 12$
$l = 1$	0.9796	0.9786	0.9780	0.9778	0.9777	0.9777	0.9778
$l = 2$	0.9209	0.9173	0.9150	0.9143	0.9138	0.9141	0.9143
$l = 3$	0.8475	0.8406	0.8361	0.8348	0.8338	0.8344	0.8349
$l = 4$	0.7852	0.7754	0.7689	0.7672	0.7658	0.7669	0.7677
$l = 5$	0.7436	0.7316	0.7236	0.7216	0.7200	0.72159	0.7228
$l = 6$	0.7170	0.7020	0.6928	0.6907	0.6889	0.6909	0.6926
$l = 7$	—	0.6891	0.6778	0.6756	0.6738	0.6761	0.6780
$l = 8$	—	—	0.6679	0.6653	0.6635	0.6661	0.6682
$l = 9$	—	—	—	0.6631	0.6609	0.6638	0.6660
$l = 10$	—	—	—	—	0.6591	0.6626	0.6648
$l = 11$	—	—	—	—	—	0.6651	0.6678
$l = 12$	—	—	—	—	—	—	0.6698

Numerical example: $\overline{D}$

	$c = 4$	$c = 5$	$c = 6$	$c = 7$	$c = 8$	$c = 9$	$c = 10$
$l = 1$	57.4139	42.4111	37.6298	37.6303	37.5259	39.0722	40.2516
$l = 2$	57.2731	42.2940	37.5278	37.5399	37.4464	39.0008	40.1824
$l = 3$	60.1768	45.0087	40.1156	40.1894	40.1384	41.8741	43.2097
$l = 4$	63.3439	48.1818	43.2285	43.3755	43.3763	45.2714	46.7414
$l = 5$	—	52.5499	47.4802	47.6944	47.7292	49.7700	51.3546
$l = 6$	—	—	51.1644	51.6063	51.6833	53.8279	55.4936
$l = 7$	—	—	—	56.5320	56.7008	58.9817	60.7385
$l = 8$	—	—	—	—	60.9422	63.4986	65.3284
$l = 9$	—	—	—	—	—	68.6834	70.6856
$l = 10$	—	—	—	—	—	—	75.3074

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Numerical complexity

- $\forall c$: calculate c roots
- $\forall (I, c)$: solve set of c equations in c unknowns
- Much numerical work
- Use results from model of Dorfman?

Static model

- Many items (N), all present from the start
- Average # tests ($E[T]$) to screen all items
- Example: individual retesting:

$$E[T] \sim \frac{N}{c} [1 + c(1 - \bar{p}^c)]$$

for $N \rightarrow \infty$

Dynamic versus static

- Group size(s)

- ▶ Static: one group size (no minimum)
- ▶ Dynamic: minimum and maximum group size (l and c)

- Performance measures

- ▶ Static: average # tests ($E[T]$)
- ▶ Dynamic: testing probability (f), mean delay of items ($\bar{D}$)

- Input parameters

- ▶ Static: # items to be screened (N)
- ▶ Dynamic: distribution of # item arrivals in a time unit ($A(z)$)

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- l_f : minimum group size that minimizes f
- c_f : maximum group size that minimizes f
- c_s : optimal static group size (minimizes $E[T]$)
- Main result: $l_f = c_f = c_s$, regardless of
 - ▶ N (static)
 - ▶ $A(z)$ (dynamic)
 - ▶ Prevalence rate

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- $l_{\overline{D}}$: minimum group size that minimizes $\overline{D}$
- $c_{\overline{D}}$: maximum group size that minimizes $\overline{D}$
- Results:
 - ▶ $c_{\overline{D}} \neq c_s$: numerical work involved
 - ▶ $l_{\overline{D}} = 1$ good heuristic: reduces complexity

Search space for $c_{\overline{D}}$ (1)

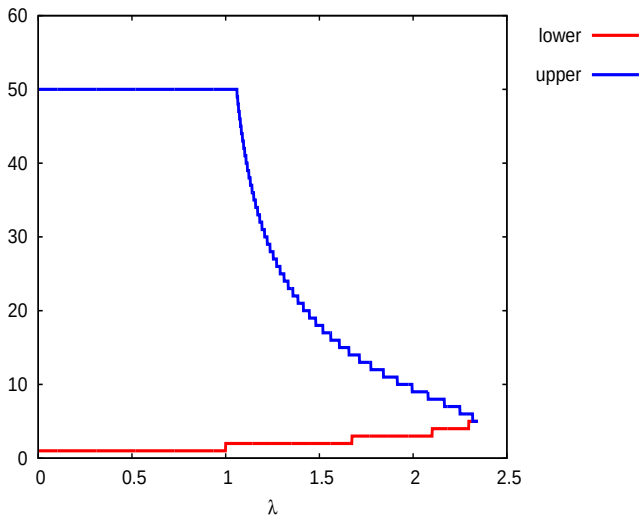
- Stability condition:

$$\rho \triangleq \lambda \frac{\mathbb{E}[S_c]}{c} < 1$$

- ▶ λ : mean arrival rate (given)
- ▶ c : group size
- ▶ S_c : # tests to screen group of c items

Search space for $c_{\overline{D}}$ (2)

- ρ smallest for $c = c_s$
- $c \uparrow$ or $c \downarrow$ (as compared to c_s) $\Rightarrow \rho \uparrow$
- c too small or too large $\Rightarrow \rho > 1$
- No numerical work



Algorithm for small λ

- Search space **largest for small λ**
- Algorithm for small λ
- Two scenarios:
 - ▶ Bursty arrivals
 - ▶ Poisson like arrivals (e.g. Poisson, Bernoulli, geometric)

Bursty arrivals

Main result: $c_{\overline{D}} = c_s$ good heuristic

Poisson like arrivals

Taylor series expansion of $\overline{D}(\lambda)$ about $\lambda = 0$:

$$\overline{D}(\lambda) = \overline{D}_0 + \lambda \overline{D}_1 + \lambda^2 \overline{D}_2 + \dots$$

Theorem

$\overline{D}_j$ takes into account the possibilities

of having $1, 2, \dots, j + 1$ arrivals

in a time unit with item arrivals

Poisson like arrivals: $\overline{D}_0$

$$\overline{D}(\lambda) = \overline{D}_0 + \lambda \overline{D}_1 + \lambda^2 \overline{D}_2 + \dots$$

- 1 item arrival
- No information about $c_{\overline{D}}$

Poisson like arrivals: $\overline{D}_1$

$$\overline{D}(\lambda) = \overline{D}_0 + \lambda \overline{D}_1 + \lambda^2 \overline{D}_2 + \dots$$

- 1 or 2 arrivals
- $c = 1$ versus $c \geq 2$
- $3/2\mathbb{E}[S_1]$ versus $\mathbb{E}[S_2]$
- If $c = 1$ best: stop; else: continue with $\overline{D}_2$

Poisson like arrivals: $\overline{D}_k$

$$\overline{D}(\lambda) = \overline{D}_0 + \dots + \lambda^k \overline{D}_k + \dots$$

- 1 or 2, or, ..., $k + 1$ arrivals
- $c = k$ versus $c \geq k$
- $E[S_k] + 1/(k + 1)E[S_1]$ versus $E[S_{k+1}]$
- If $c = k$ best: stop; else: continue with $\overline{D}_{k+1}$

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- Queueing model to include dynamic nature of item arrivals
 - ▶ Batch/bulk service
 - ▶ Screening policy captured in service times
 - ▶ Testing probability f and mean delay $\overline{D}$
 - ▶ Requires (much) numerical work
- When are static results useful in dynamic context?

```
if optimization criterion == f
    optimal l = optimal c = c_s;
if optimization criterion == delay
    optimal l = 1 (heuristic);
    if small arrival_rate
        if bursty arrivals
            optimal c = c_s (heuristic);
        if Poisson arrivals
            efficient algorithm for optimal c
    else
        determine search space
        calculate delay for all group sizes from
        the search space and select the group size
        that produces smallest delay
```

Future work

- $c_{\overline{D}} \leq c_s$?
- $c_{\overline{D}}$ non-decreasing function of λ ?
- Accurate closed-form approximation for $\overline{D}$ for medium λ
- Same conclusions in case of False Positives and False Negatives?

Questions?

